



## GRADED EXERCISE SHEET

### Definition 1.

- An undirected graph is a tuple  $(V, E)$  where  $V$  is any set and  $E \subseteq V \times V$  such that for all  $v, w \in V$  it holds that  $(v, w) \in E \Rightarrow (w, v) \in E$ .
- A graph isomorphism between graphs  $(V, E), (V', E')$  is a bijection  $f : V \rightarrow V'$  such that  $(v, w) \in E \Leftrightarrow (f(v), f(w)) \in E'$ .
- A countable graph  $(V, E)$  is called **computable** if it is either finite or isomorphic to some graph  $(\mathbb{N}, E')$  such that  $E' \subseteq \mathbb{N}^2$  is a computable relation.

There is an obvious numbering  $\Gamma$  of the isomorphism classes of countable graph:

The domain of the numbering is

$$\text{dom } \Gamma = \{i \mid i \in \text{Tot} \wedge \text{im}(\varphi_i) \subseteq \{0, 1\} \wedge \varphi_i(n, m) = 1 \Leftrightarrow \varphi_i(m, n) = 1\}$$

and the numbering is defined by

$$\forall i : \Gamma_i = (\mathbb{N}, E_{\varphi_i})$$

where  $E_{\varphi_i}$  is the relation defined by the characteristic function  $\varphi_i$ .

### Exercise 1. (Domain of $\Gamma$ )

Place  $\text{dom } \Gamma$  in the arithmetical hierarchy.

### Definition 2.

- A path in some graph  $(V, E)$  is a sequence of edges  $e_1, \dots, e_n \in E$  such that for  $1 \leq i < n$  if  $e_i = (v, w)$  then  $e_{i+1} = (w, v')$ .
- A graph is connected if for any two vertices there is a path starting at one and ending at the other.
- The distance of two vertices  $v, w$  of some graph  $(V, E)$  is the length of a shortest path between  $v$  and  $w$ . We denote it by  $d(v, w)$ .

### Exercise 2. (Connectedness)

Place the set

$$\text{Con} = \{i \mid i \in \text{dom } \Gamma \wedge \Gamma_i \text{ is connected}\}$$

in the arithmetical hierarchy.

**Exercise 3. (Infinite line)**

Place the set

$$Lin = \{i \mid i \in \text{dom } \Gamma \wedge \Gamma_i \text{ is a bi-infinite line}\}$$

in the arithmetical hierarchy.

**Exercise 4. (Uncomputable distance)**

- a) Show that on any computable graph the distance function  $d : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  is upper computable, i.e. there is a computable function  $\Phi : \mathbb{N} \times \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  such that for fixed  $n, m$  the sequence  $(\Phi(n, m, i))_{i \in \mathbb{N}}$  is monotonically decreasing and converges to  $d(n, m)$ .
- b) Give an example of a computable graph such that the distance function is not computable.

**Exercise 5. (New Horizons)**

Find a graph theoretical problem  $A \subseteq \text{dom } \Gamma$  and place it in the arithmetical hierarchy. It should be strictly higher than  $\text{dom } \Gamma$ .