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EXERCISE SHEET NO. 4 FOR COMPUTABILITY IN MATHEMATICS

## Exercise 1. (Reductions of c.e. sets - 3 Points)

Prove the following
a) If $A \leq_{m} B$ and $B$ is c.e., then $A$ is c.e.
b) $A$ is c.e. in $B$ if and only if $A$ is c.e. in $\bar{B}$.
c) If $A$ is c.e. in $B$ and $B \leq_{T} C$ then $A$ is c.e. in $C$.

## Exercise 2. (Cylinders - 3 Points)

A set $A$ is a cylinder if $\forall B: B \leq_{m} A \Rightarrow B \leq_{1} A$. Show that this is equivalent to $A \equiv_{1} C \times \mathbb{N}$ for some set $C$.

## Exercise 3. (Lattice structure of Turing degrees - 4 Points)

For sets $A, B \subseteq \mathbb{N}$ we define $A \oplus B=\{2 x \mid x \in A\} \cup\{2 x+1 \mid x \in B\}$.
a) Let $A$ be some set and $B=A \oplus \bar{A}$. Prove $B \leq_{1} \bar{B}$.
b) Show that $A, B \leq_{T} A \oplus B$.
c) Show that if $A, B \leq_{T} C$, then $A \oplus B \leq_{T} C$.
d) Conclude that in the set of Turing degrees, any two elements have a common least upper bound.

## Exercise 4. (Cardinality of Turing degrees - 4 Points)

Prove the following:
a) Each Turing degree contains exactly $\aleph_{0}$ many sets.
b) For each Turing degree $\mathbf{a}$ the set of Turing degrees $\mathbf{b}$ with $\mathbf{b}<\mathbf{a}$ is countable.
c) There are at least $2^{\aleph_{0}}$ different Turing degrees.
d) There are exactly $2^{\aleph_{0}}$ different Turing degrees.

Consider the following inductive definition of two sequences of partial functions $f_{i}, g_{i}: \mathbb{N} \rightarrow\{0,1\}$ :
$\mathrm{i}=0 f_{0}=g_{0}$ are undefined everywhere.
$i=2 e$ Choose the smallest $x \in \mathbb{N}$ such that $g_{i-1}(x)$ is undefined. If there is a finite extension $f_{i+1}$ of $f_{i}$ such that $\varphi_{i}^{f_{i+1}}(x)$ halts and has a value in $\{0,1\}$, then fix this $f_{i+1}$ and let $g_{i+1}$ be the extension of $g_{i}$ by $g_{i}(x)=1-\varphi_{i}^{f_{i+1}}(x)$. If there is not such a finite extension, define $f_{i+1}=f_{i}$ and $g_{i+1}$ to be the extension of $g_{i}$ by $g_{i+1}(x)=0$.
$i=2 e+1$ Proceed as in the case $i=2 e$ with exchanged roles of $f$ and $g$.
Note: $\varphi_{e}^{f}$ refers to the OTM with number $e$ which has the partial characteristic function $f$ on its tape.

## Exercise 5. (Kleene-Post-Theorem - 4 Points)

a) Prove that for all $e$ it holds that $\varphi_{e}^{f_{2 e}} \neq g_{2 e}$ and $\varphi_{e}^{g_{2 e+1}} \neq f_{2 e+1}$.
b) Prove that there exist incomparable Turing degrees $\mathbf{a}$ and $\mathbf{b}$, i.e. $\mathbf{a} \not \leq_{T} \mathbf{b}$ and $\mathbf{b} \not \leq_{T} \mathbf{a}$. This is known as the Kleene-Post-Theorem.

