

# EXERCISE SHEET NO. 4 FOR COMPUTABILITY IN MATHEMATICS

## Exercise 1. (Reductions of c.e. sets - 3 Points)

Prove the following

- a) If  $A \leq_m B$  and B is c.e., then A is c.e.
- b) A is c.e. in B if and only if A is c.e. in  $\overline{B}$ .
- c) If A is c.e. in B and  $B \leq_T C$  then A is c.e. in C.

## Exercise 2. (Cylinders - 3 Points)

A set A is a **cylinder** if  $\forall B : B \leq_m A \Rightarrow B \leq_1 A$ . Show that this is equivalent to  $A \equiv_1 C \times \mathbb{N}$  for some set C.

### Exercise 3. (Lattice structure of Turing degrees - 4 Points)

For sets  $A, B \subseteq \mathbb{N}$  we define  $A \oplus B = \{2x \mid x \in A\} \cup \{2x+1 \mid x \in B\}$ .

- a) Let A be some set and  $B = A \oplus \overline{A}$ . Prove  $B \leq_1 \overline{B}$ .
- b) Show that  $A, B \leq_T A \oplus B$ .
- c) Show that if  $A, B \leq_T C$ , then  $A \oplus B \leq_T C$ .
- d) Conclude that in the set of Turing degrees, any two elements have a common least upper bound.

#### Exercise 4. (Cardinality of Turing degrees - 4 Points)

Prove the following:

- a) Each Turing degree contains exactly  $\aleph_0$  many sets.
- b) For each Turing degree  $\mathbf{a}$  the set of Turing degrees  $\mathbf{b}$  with  $\mathbf{b} < \mathbf{a}$  is countable.
- c) There are at least  $2^{\aleph_0}$  different Turing degrees.
- d) There are exactly  $2^{\aleph_0}$  different Turing degrees.

Consider the following inductive definition of two sequences of partial functions  $f_i, g_i : \mathbb{N} \to \{0, 1\}$ :

i=0  $f_0 = g_0$  are undefined everywhere.

- i = 2e Choose the smallest  $x \in \mathbb{N}$  such that  $g_{i-1}(x)$  is undefined. If there is a finite extension  $f_{i+1}$  of  $f_i$  such that  $\varphi_i^{f_{i+1}}(x)$  halts and has a value in  $\{0, 1\}$ , then fix this  $f_{i+1}$  and let  $g_{i+1}$  be the extension of  $g_i$  by  $g_i(x) = 1 \varphi_i^{f_{i+1}}(x)$ . If there is not such a finite extension, define  $f_{i+1} = f_i$  and  $g_{i+1}$  to be the extension of  $g_i$  by  $g_i(x) = 0$ .
- i = 2e + 1 Proceed as in the case i = 2e with exchanged roles of f and g.

Note:  $\varphi_e^f$  refers to the OTM with number e which has the partial characteristic function f on its tape.

## Exercise 5. (Kleene-Post-Theorem - 4 Points)

- a) Prove that for all e it holds that  $\varphi_e^{f_{2e}} \neq g_{2e}$  and  $\varphi_e^{g_{2e+1}} \neq f_{2e+1}$ .
- b) Prove that there exist incomparable Turing degrees **a** and **b**, i.e.  $\mathbf{a} \not\leq_T \mathbf{b}$  and  $\mathbf{b} \not\leq_T \mathbf{a}$ . This is known as the Kleene-Post-Theorem.