



EXERCISE SHEET NO. 1 FOR COMPUTABILITY IN MATHEMATICS

Exercise 1. (Turing Machine manipulations - 4 Points)

- Implement Euclid's algorithm for the gcd on a Turing machine.
- Give a machine that, on input a word, rewrites this word moved one left on the tape, and one that moves it to the right.
Use this to describe an "insertion" move for Turing machines: expressiveness of Turing machines is not increased by allowing them to insert a tape symbol at a desired position, instead of rewriting over what there is.

Exercise 2. (Composition is computable - 4 Points)

- Show that if $f : \subseteq A^* \rightarrow A^*$ is computable, then it can be computed by a machine M , with initial state q_0 and final state q_f , such that for any w in $\text{dom}(f)$, $q_0 w \rightarrow_M^* q_f f(w)$.
(In other words: the machine computes f and then positions correctly its reading head at the beginning of the output.)
Prove that the composition of computable functions is computable.
- Strengthen the previous point in a uniform way: show that the map which associates to a pair of codes for Turing machines the code of a machine that computes their composition is computable.
This could simply be phrased as: "show that the composition function $\circ : \mathcal{PC} \times \mathcal{PC} \rightarrow \mathcal{PC}$ is computable".

Exercise 3. (Cantor Pairing functions - 3 Points)

Cantor's Pairing function is the map $\pi : (x, y) \mapsto (x^2 + 2xy + 3x + y^2 + y)/2$.

- Justify that, restricted to the set A_u , $u \in \mathbb{N}$, given by $A_u = \{(x, y) \in \mathbb{N}^2, x + y = u\}$, π realizes a bijection onto $\{(u^2 + u)/2, \dots, (u^2 + 3u)/2\}$
- Justify that if $x + y \neq z + t$, $\pi(x, y) \neq \pi(z, t)$.
- Show that π is a bijection.

Exercise 4. (Left Moves in a TM - 4 Points)

Prove that it is decidable whether a Turing machine ever in its computation for a given input moves the tape pointer to the left. *Hint: Understand the case of input having size 0 first.*